



CS 5/7320
Artificial Intelligence

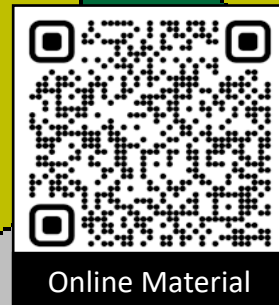
Constraint Satisfaction Problems

AIMA Chapter 6

Slides by Michael Hahsler
based on Slides by Svetlana Lazepnik
with figures from the AIMA textbook

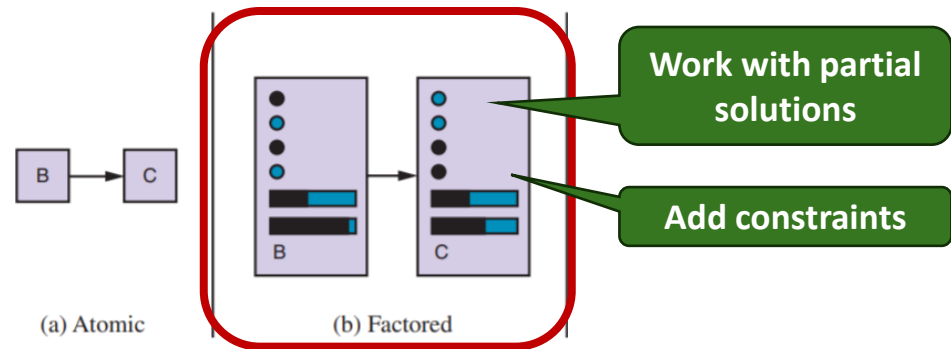


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Constraint Satisfaction Problems (CSPs)



Definition:

- **State** is defined by a factored state representation:
 - A set of variables X_i called fluents.
- **Partial Solution**: Each variable can have a value from domain D_i or be **unassigned**.
- **Constraints** are a set of rules specifying allowable combinations of values for subsets of the variables.

E.g., $X_1 = 3$, $X_1 \neq X_7$ or $X_2 > X_9 + 3$

- **Solution**: a state that is a
 - Consistent assignment**: satisfies all constraints.
 - Complete assignment**: assigns value to each variable.



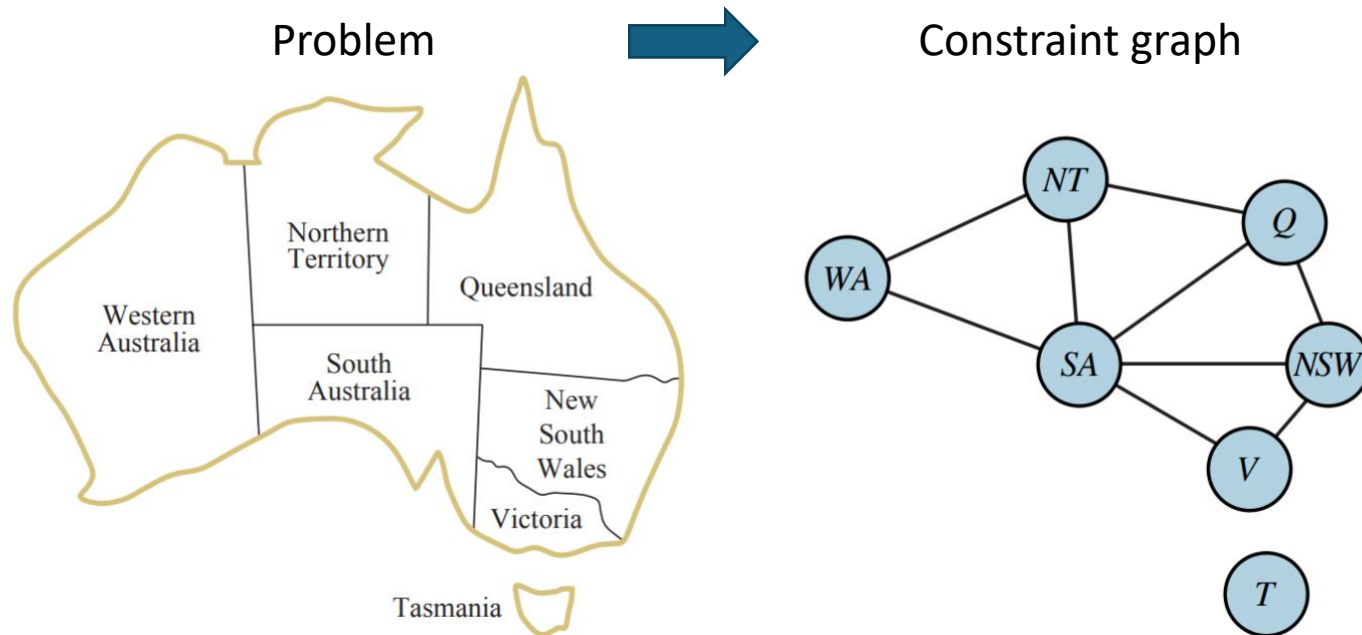
Comparison to Other Methods

	Generic Tree Search	Local Search	CSP
State representation	Atomic states Variables are only used to create human readable labels or calculate heuristics.	Factored representation to find local moves.	Factored
Assignment	Always complete	Always complete	Partial assignment during search
Constraints	Constraints are implicit in the search problem (initial + goal state + transition function).	Constraints are represented by the objective function and may not be met.	Enforces explicit constraints .

+ General-purpose solvers for CSP with more power than standard search algorithms exist.



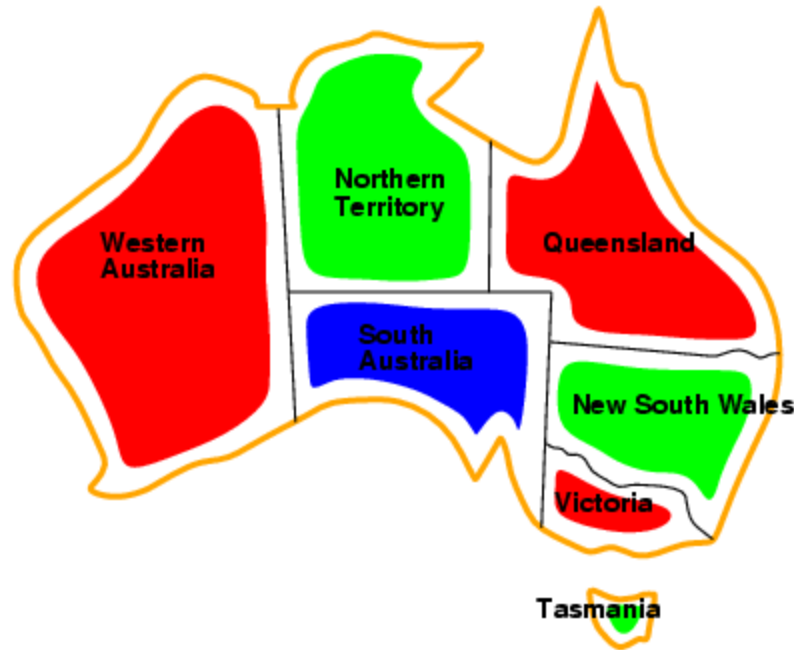
Example: Map Coloring (Graph coloring)



- **Variables representing state:** WA, NT, Q, NSW, V, SA, T
- **Variable Domains:** {red, green, blue}
- **Constraints:** adjacent regions must have different colors
e.g.,
$$WA \neq NT \Leftrightarrow (WA, NT) \in \{(\text{red}, \text{green}), (\text{red}, \text{blue}), (\text{green}, \text{red}), (\text{green}, \text{blue}), (\text{blue}, \text{red}), (\text{blue}, \text{green})\}$$



Example: Map Coloring



Solutions are *complete* and *consistent* assignments, e.g.,

WA = red, NT = green, Q = red, NSW = green,
V = red, SA = blue, T = green



Example: N-Queens

- **Variables:** X_{ij} for $i, j \in \{1, 2, \dots, N\}$
- **Domains:** $\{0, 1\}$ # Queen: no/yes

- **Constraints:**

$$\sum_{i,j} X_{ij} = N$$

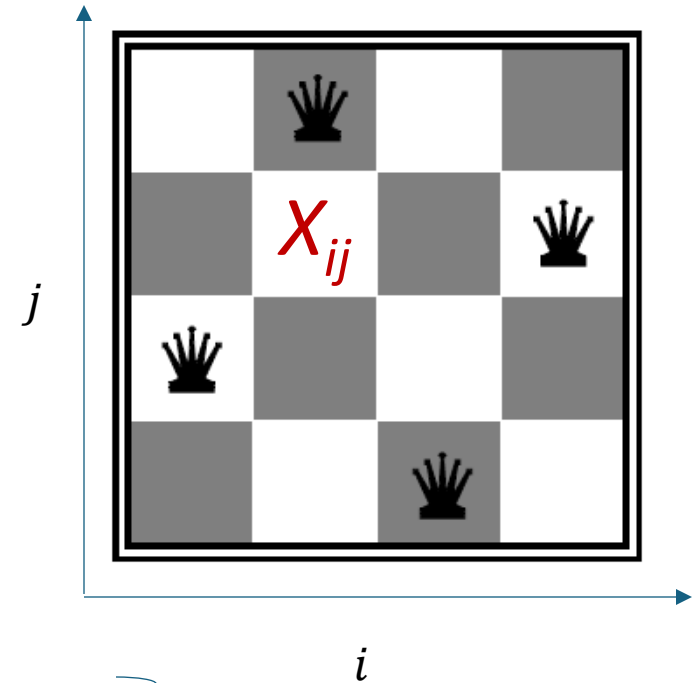
$(X_{ij}, X_{ik}) \in \{(0, 0), (0, 1), (1, 0)\}$ # cannot be in same col.

$(X_{ij}, X_{kj}) \in \{(0, 0), (0, 1), (1, 0)\}$ # cannot be in same row.

$(X_{ij}, X_{i+k, j+k}) \in \{(0, 0), (0, 1), (1, 0)\}$ # cannot be diagonal

$(X_{ij}, X_{i+k, j-k}) \in \{(0, 0), (0, 1), (1, 0)\}$ # cannot be diagonal

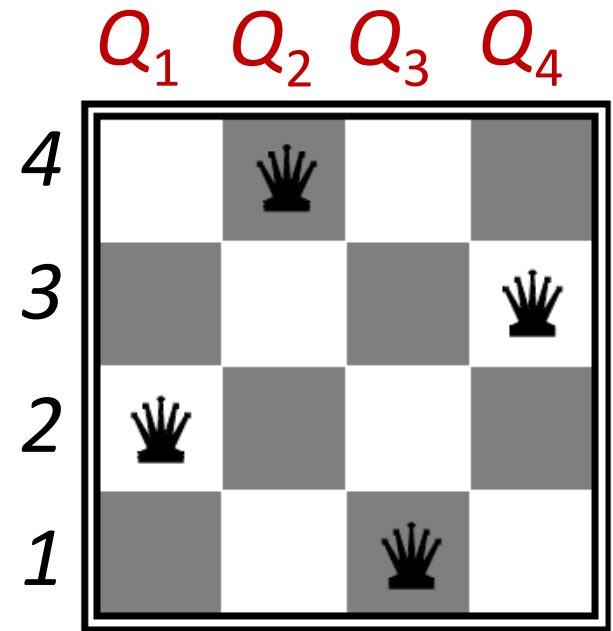
for $i, j, k \in \{1, 2, \dots, N\}$





N-Queens: Alternative Formulation

- **Variables:** Q_1, Q_2, \dots, Q_N
- **Domains:** $\{1, 2, \dots, N\}$ # row for each col.
- **Constraints:**
 $\forall i, j$ non-threatening (Q_i, Q_j)



Example:

$Q_1 = 2, Q_2 = 4, Q_3 = 1, Q_4 = 3$

Example: Sudoku

- **Variables:** X_{ij}
- **Domains:** $\{1, 2, \dots, 9\}$
- **Constraints:**
 - Alldiff(X_{ij} in the same *unit*)
 - Alldiff(X_{ij} in the same *row*)
 - Alldiff(X_{ij} in the same *column*)

					8			4
	8	4		1	6			
			5			1		
1		3	8			9		
6		8		X_{ij}		4		3
		2			9	5		1
		7			2			
			7	8		2	6	
2			3					



Some Popular Types of CSPs

- **Boolean Satisfiability Problem (SAT)**

Find variable assignments that make a Boolean expression (often expressed in conjunctive normal form) evaluate as true.

$$(x_1 \vee \neg x_2) \wedge (\neg x_1 \vee x_2 \vee x_3) \wedge \neg x_1 = \text{True}$$

- **Integer Programming**

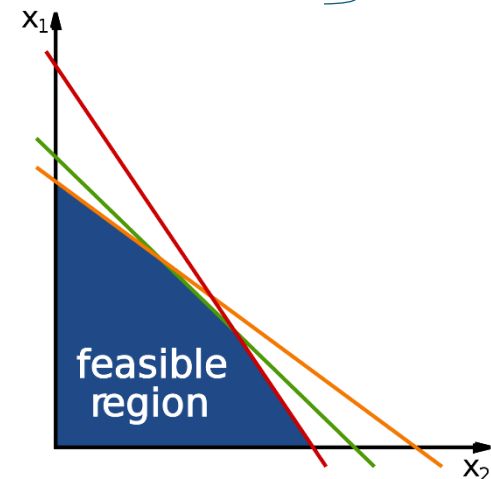
Variables are restricted to integers. Find a feasible solution that satisfies all constraints. The traveling salesman problem can be expressed as an integer program.

- **Linear Programming**

Variables are continuous and constraints are linear (in)equalities.

Find a feasible solution using, e.g., the simplex algorithm.

NP-complete





Real-world CSPs

- **Assignment problems**

e.g., who teaches what class for a fixed schedule. A teacher cannot be in two classes at the same time!

- **Timetable problems**

e.g., which class is offered when and where? No two classes in the same room at the same time.

- **Scheduling** in transportation and production (e.g., order of production steps).

- Many problems can naturally also be formulated as CSPs.

- More examples of CSPs: <http://www.csplib.org/>



Constraint Satisfaction Problems and Tree Search



Formulation of a CSP as a Search Problem

State:

- Values assigned so far

Initial state:

- The empty assignment $\{ \}$ (all variables are unassigned)

Transition function (Successor function):

- Choose an unassigned variable and assign it a value that does not violate any constraints
- Fail if no legal assignment is found

Goal state:

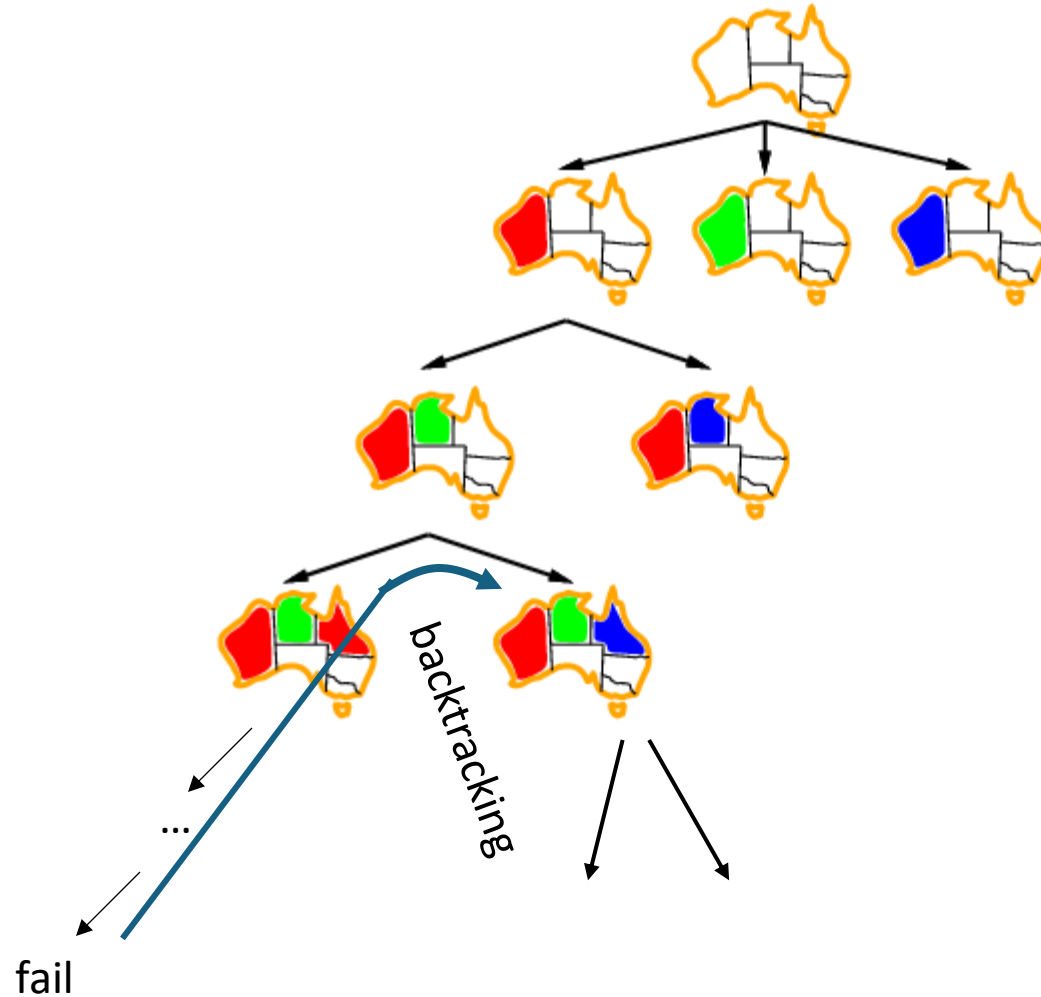
- Any complete and consistent assignment.

Note: Path cost is not important.

Backtracking Search

- Start with all variables unassigned.
- Build a search tree that assigns a value to one variable per level.
 - Tree depth: number of variables n
 - Number of leaves: $O(d^n)$ where d is the number of values per variable
- Depth-first search for CSPs with single-variable assignments is typically done with **backtracking search**.

Example: Backtracking Search (DFS)



Backtracking Search Algorithm

```
function RECURSIVE-BACKTRACKING(assignment, csp)  
  if assignment is complete then return assignment  
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)  
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp)  
    if value is consistent with assignment given CONSTRAINTS[csp]  
      add {var = value} to assignment  
      result ← RECURSIVE-BACKTRACKING(assignment, csp)  
      if result ≠ failure then return result  
      remove {var = value} from assignment  
  return failure
```

Call: Recursive-Backtracking({}, *csp*)

Improving backtracking efficiency:

- Which variable should be assigned next?
- In what order should its values be tried?
- Can we detect inevitable failure early?

Similar to move ordering in games.

Tree pruning (like in alpha-beta search)



Local Search for Constraint Satisfaction Problems

Minimize violated constraints



Local Search for CSPs

CSP algorithms

- Allow incomplete states.
- States must satisfy all constraints.

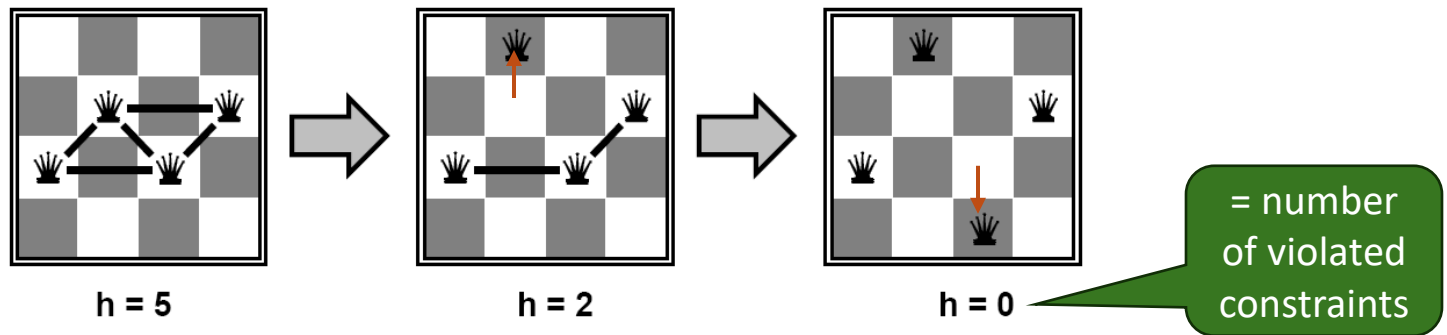
VS.

Local Search

- Only complete states
- Allows states with unsatisfied constraints.

Local search can attempt to reduce unsatisfied constraints using the **min-conflicts** heuristic:

1. Select a variable that violates a constraint (produces a conflict).
2. Choose a new value that violates fewer constraints.
3. Repeat till all constraints are met (or a local optimum is reached).



Simulated Annealing is often very effective for CSPs. Especially for very large problems where an imperfect solution is acceptable.



What You Should Know

- CSPs are a special type of search problem:
 - States are **factored** and defined by a set of variables and value assignments
 - The goal is defined by a set of **constraints** on the variables.
 - Incomplete assignments are used to create a complete assignment piece by piece.
 - The goal test is defined by
 - **Consistency** with constraints
 - **Completeness** of assignment
- **Many problems can be formulated as a CSP**, and problems where the constraints are very restrictive on the solution space may be easier to solve as CSPs (e.g., scheduling problems and puzzles).
- Effective off-the-shelf solvers exist. They typically use:
 - **Depth-first search**: successor states are generated variable-by-variable by adding a consistent value assignment to single unassigned variables.
 - **Local search** can be used as an effective heuristic. It searches the space of all complete assignments for consistent assignments = **min-conflicts heuristic**.